

Length-Based Anchor-Free Localization in a Fully Covered Sensor Network

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First International Conference on COMmunication Systems
and NETworkS (COMSNETS), January 5-10, 2009,
Bangalore, India

Outline

- 1 **Introduction**
- 2 **Basic Model and Problem Statement**
- 3 **Localization under the Basic Model**
- 4 **Restricted Model**
 - Localization under the Restricted Model
 - Correctness and Complexity
 - Simulation Studies
- 5 **Conclusion**

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Introduction

Applications of WSN

WSN has wide-ranging applications

- habitat monitoring
- battlefield surveillance
- monitoring forest-fire
- disaster management etc.

Introduction

Localization

Why is the localization a major issue in WSN?

- Sensors may be **deployed randomly**,
- positions of the sensors are completely unknown,
- positions are extremely important in **routing, security** etc.

Introduction

Localization

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- positions of the sensors are completely unknown,
- positions are extremely important in **routing, security** etc.

Definition

Finding the positions of sensors is known as **localization** in WSN.

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Basic Model

- Field of interest is **convex**.
- Sensors are **identical**.
- **Sensing range** $\rightarrow s$. **Communication range** $\rightarrow r$.
- **No barrier** present in the WSN field.
- Sensors are in **general position**.
- Distances are measured **accurately**.
- For (s_i, s_j) ,
if $dist(s_i, s_j) \leq r$, the value of $dist(s_i, s_j)$ is **known**.
If $dist(s_i, s_j) > r$, the value of $dist(s_i, s_j)$ is **unknown**.

Basic Model

Graph Model of WSN

Let WSN contain n sensors s_1, s_2, \dots, s_n . We construct an undirected edge-weighted graph $G = (V, E, w)$ as follows :

- Each vertex in V represents a sensor.
- $(s_i, s_j) \in E$ iff the distance between s_i and s_j is known.
- $\forall (s_i, s_j) \in E, w(s_i, s_j) = \text{dist}(s_i, s_j)$.

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Goal

To localize the sensors without using any anchor.

Basic Model

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Goal

To **localize** the sensors **without** using any **anchor**.

Definition

A **realization** of $G = (V, E, w)$ in \mathbb{R}^m (Euclidean space of dimension m) is an assignment of coordinates (x_1, \dots, x_m) to the vertices so that weight of an edge represents the Euclidean distance between the points incident on the edge.

Problem Statement

We are interested only in \mathbb{R}^2 .

Problem

Given an *edge-weighted undirected graph* $G = (V, E, w)$, our objective is to find *all possible realizations* of the graph G in \mathbb{R}^2 such that,

$$\text{dist}(s_i, s_j) > r, \quad \forall (s_i, s_j) \notin E.$$

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Sensor realization

A *solution to this problem* will be referred to as a *sensor realization* of G .

The sensor realization of the graph G may be denoted by $\mathcal{R} = (G, \Phi)$, where $\Phi : V \rightarrow \mathbb{R}^2$ and $\Phi(v)$ represents the corresponding *realization of the vertex*, $v \in V$.

Problem Statement

Is there any Sensor realization of the WSN at all?

- **At least one sensor realization exists**, since the distance information is coming from an actual deployment of sensors

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Localization under the Basic Model

In general

- WSN may have **multiple Sensor realizations**.
- Actually, any sensor realization may be flipped, rotated and/or translated (shifting origin), (**like any rigid body**), to get another sensor realization.
- We are interested only in realizations which are **structurally different**.

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- We are interested only in realizations which are **structurally different**.

Our main goal is to find only:

- A set of **sensor realizations** of a graph from which any other sensor realization may be obtained with the operation **shifting**, **flipping** and/or **rotation** suitably.

Localization under the Basic Model

Rigidity, Flipping, Rotation and Shifting

- From here onwards **flip**, **rotation** and **shift** in a sensor realization of a WSN graph, we mean a **part of the sensor realization** is flipped, rotated or shifted, giving us a new sensor realization while rest of the realization remains fixed.

Localization under the Basic Model

Rigidity, Flipping, Rotation and Shifting

Definition

An edge-weighted graph $G = (V, E, w)$ is called **rigid**, if the **distance** between any pair of sensors remains same in **every sensor realization**.

Localization under the Basic Model

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An edge-weighted graph $G = (V, E, w)$ is called **rigid**, if the **distance** between any pair of sensors remains same in **every sensor realization**.

Theorem

If two rigid bodies B_1 and B_2 , in a sensor realization of a graph $G = (V, E, w)$ share three or more common vertices, $B_1 \cup B_2$ (alongwith all edges between them) forms a rigid body.

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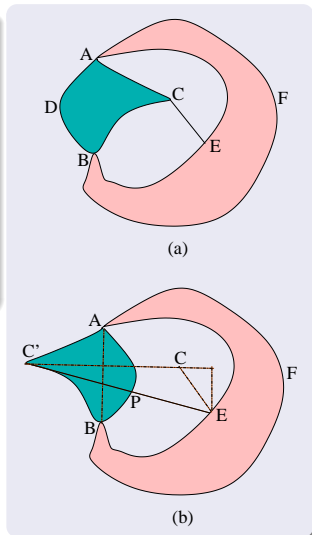
If two rigid graphs share exactly one vertex, then one of them may be rotated, around the common vertex, keeping the other fixed. Such a vertex will be called a **joint**.

Localization under the Basic Model

Rigidity, Flipping, Rotation and Shifting

Definition

If two rigid graphs share exactly two vertices, rotation about these vertices is no longer possible, but one may be flipped, about the line joining the two common vertices, keeping the other fixed. These two vertices will be called a **flip joint** or simply **flip** and each vertex as a **flip vertex**.



Localization under the Basic Model

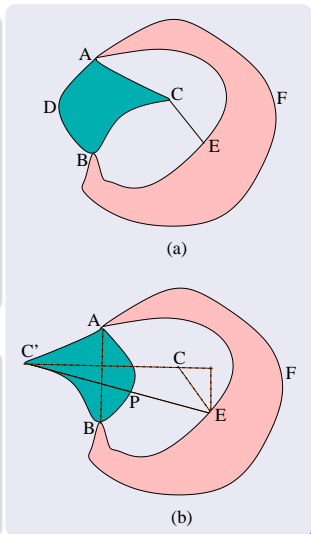
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Theorem

If two rigid subgraphs, B_1 and B_2 , of a graph G share two common vertices and there is an edge connecting a vertex in B_1 to another vertex in B_2 , $B_1 \cup B_2$ is rigid.



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Restricted Model

Characteristics

- 1 The field \mathcal{F} is **sensing covered** i.e., every point in \mathcal{F} is in the sensing zone of at least one sensor.

Restricted Model

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In VigilNet surveillance system

He et al. (*ACM Transactions on Sensor Networks*, 2006) uses **full coverage** of the field in VigilNet surveillance system. This may be required by many other application to guarantee a quality of service.

Restricted Model

Characteristics

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- 2 $s \leq \frac{r}{2}$.

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- 2 $s \leq \frac{r}{2}$.

Is $s \leq \frac{r}{2}$ practical ?

Restricted Model

Characteristics

- 1 The field \mathcal{F} is **sensing covered** i.e., every point in \mathcal{F} is in the sensing zone of at least one sensor.
- 2 $s \leq \frac{r}{2}$.

Is $s \leq \frac{r}{2}$ practical ?

Sensor mote	Commu range, r	Sensing range, s	Range ratio, $\frac{s}{r}$
Mica2 (in VigilNet surveillance system)	10m	3m	$\frac{1}{3}$
Sensoria sGate (WSN platform by Sensoria Corporation)	500m	< 100m	< $\frac{1}{5}$
with Seismic sensors	$\geq 100m$	50m	$\leq \frac{1}{2}$

Outline

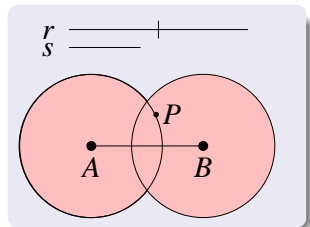
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Localization under the Restricted Model

Full Coverage verification and Unique solution

Theorem

A WSN under the restricted model is fully covered if and only if for every sensor, any point on the boundary of its sensing zone is covered by at least another sensor.

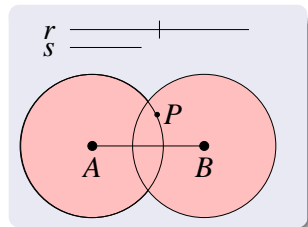


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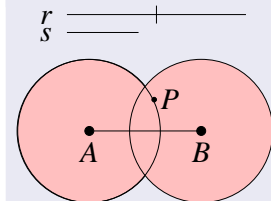
*The edge-weighted graph $G = (V, E, w)$ under the above restricted model is **rigid**.*

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Theorem

*The edge-weighted graph $G = (V, E, w)$ under the above restricted model is **rigid**.*

*If $s > \frac{r}{2}$, then the graph **may not be rigid** even if all points are covered by some sensors.*

Localization under the Restricted Model

Localization Algorithm

We propose a **centralized localization** technique.

- Its worst case complexity time is $O(|E|)$, E is the edge set of G .
- Under random **uniform deployment** of nodes, the **expected value** of $|E|$ may be represented as $O(\frac{n^2}{R})$, where R is the field area.
- The **expected time complexity** of the proposed algorithm is **significantly lower than $O(n^2)$** , in particular for **large scale WSN**.

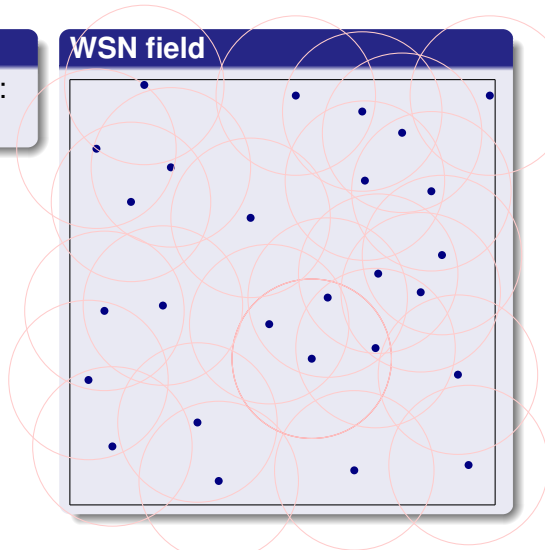
Localization under the Restricted Model

Localization Algorithm

Data structures:

- `located, locateIn2s :`
`stack;`

WSN field



Localization under the Restricted Model

Localization Algorithm

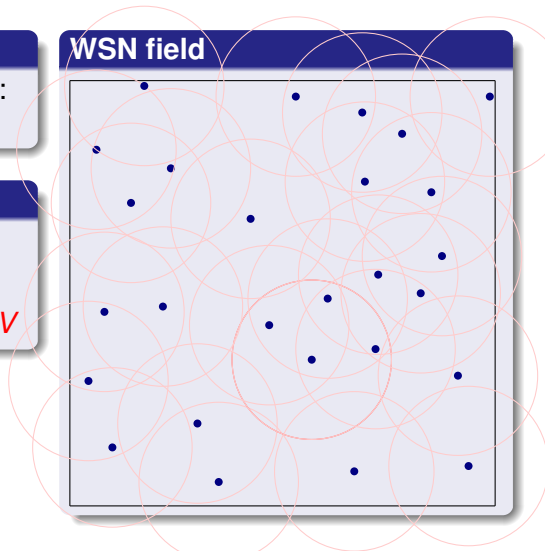
Data structures:

- `located, locateIn2s : stack;`

Initialization

- Set `located` $\leftarrow \emptyset$,
`locateIn2s` $\leftarrow \emptyset$
- Set $\Phi(v) \leftarrow \text{null}, \forall v \in V$

WSN field



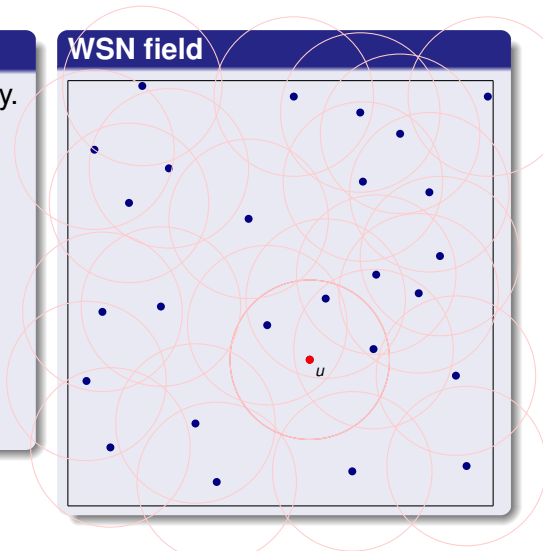
Localization under the Restricted Model

Localization Algorithm

Fixing coordinate system

- **Consider** $u \in V$ randomly.
- **Fix** $\phi(u) \leftarrow (0, 0)$

WSN field



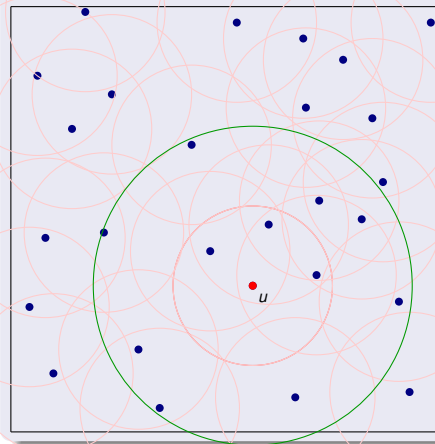
Localization under the Restricted Model

Localization Algorithm

Fixing coordinate system

- Consider $u \in V$ randomly.
- Fix $\phi(u) \leftarrow (0, 0)$
- Let $v \in C(u, 2s) - \{u\}$

WSN field



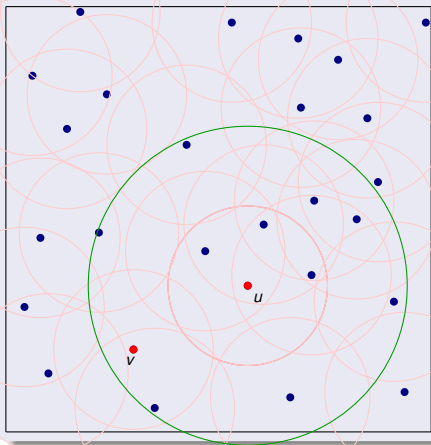
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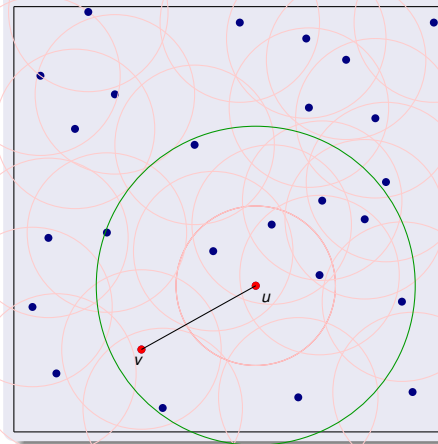
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- Fix $\phi(v) \leftarrow (w(u, v), 0)$

WSN field



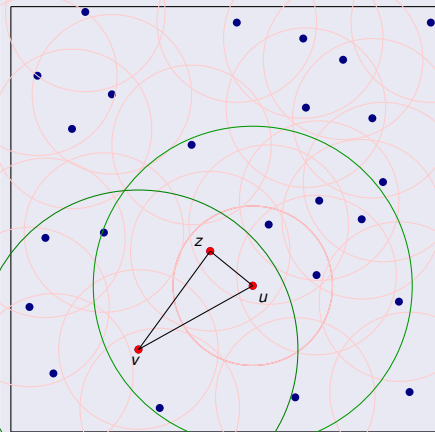
Localization under the Restricted Model

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- Fix $\phi(u) \leftarrow (0, 0)$
- Let $v \in C(u, 2s) - \{u\}$
- Fix $\phi(v) \leftarrow (w(u, v), 0)$
- Let $z \in C(u, 2s) \cap C(v, 2s) - \{u, v\}$
- Fix $\phi(z)$ satisfying $dist(u, z) = w(u, z)$ and $dist(v, z) = w(v, z)$.

WSN field



Localization under the Restricted Model

Localization Algorithm

Find $z(x_3, y_3)$ using :

$$x_3 = x_1 + \frac{x_2 - x_1}{d_3} x'_3 - \frac{y_2 - y_1}{d_3} y'_3,$$

$$y_3 = y_1 + \frac{y_2 - y_1}{d_3} x'_3 + \frac{x_2 - x_1}{d_3} y'_3$$

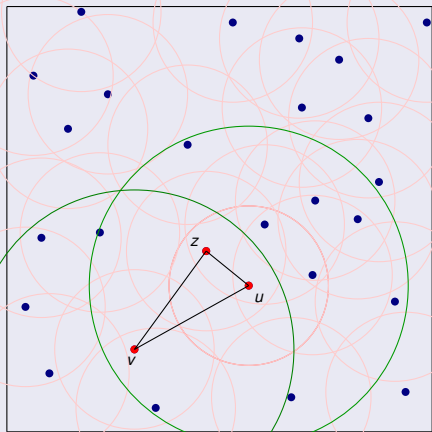
$$x'_3 = \frac{d_3^2 + d_2^2 - d_1^2}{2d_3},$$

$$y'_3 = \pm \sqrt{d_2^2 - x_3'^2}.$$

assuming

- the positions of u and v as (x_1, y_1) and (x_2, y_2) and
- $d_1 = w(v, z)$, $d_2 = w(z, u)$ and $d_3 = w(u, v)$.

WSN field



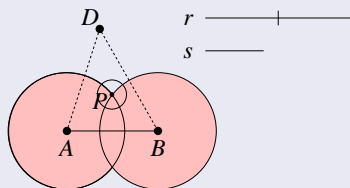
Localization under the Restricted Model

Localization Algorithm

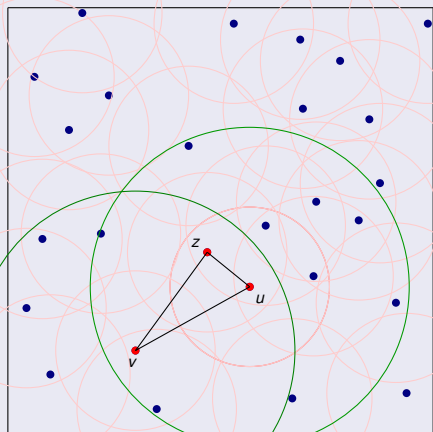
Theorem

Let $\Phi(u) = A$ and $\Phi(v) = B$. If P is an intersection point of the boundaries of $C(A, s)$ and $C(B, s)$, then P is covered by another sensor, say, D .

Existence of $D = \Phi(z)$



WSN field



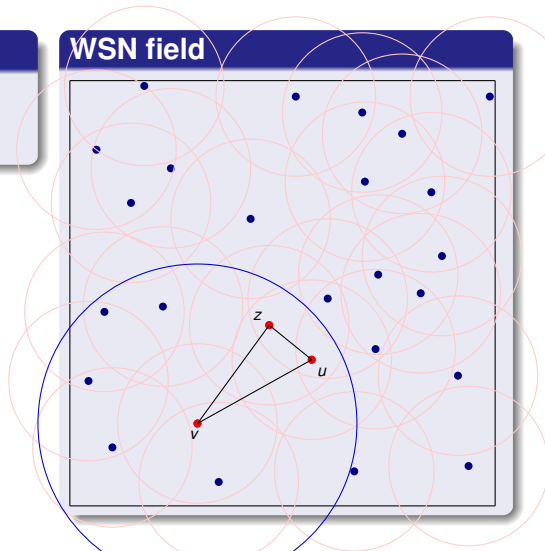
Localization under the Restricted Model

Localization Algorithm

Finding Positions

Push u and z to located
Set $s_1 \leftarrow v$

WSN field



Localization under the Restricted Model

Localization Algorithm

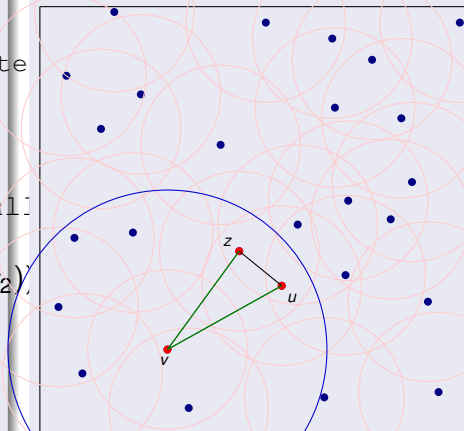
Finding Positions

```

while(located  $\neq \emptyset$ )
  Push Top(located) to locate
  while(locateIn2s  $\neq \emptyset$ )
     $s_2 \leftarrow \text{Pop}(\text{locateIn2s})$ 
    forall( $s_3 \in C(s_1, 2s) \cap$ 
       $C(s_2, 2s)$  and  $\Phi(s_3) = \text{null}$ )
       $\Phi(s_3) \leftarrow$ 
        GETPOS( $s_3, s_1, s_2, \Phi(s_2)$ )
      Push  $s_3$  to located
      and to locateIn2s
    end for
  end while
   $s_1 \leftarrow \text{Pop}(\text{located})$ 

```

WSN field



Localization under the Restricted Model

Localization Algorithm

func GETPOS(s_3, s_1, s_2, p_2)

*/** $\Phi(s_1)$ is fixed and s_2 is assigned a position, p_2 . If possible return correct position of s_3 . Otherwise, return null.**/*

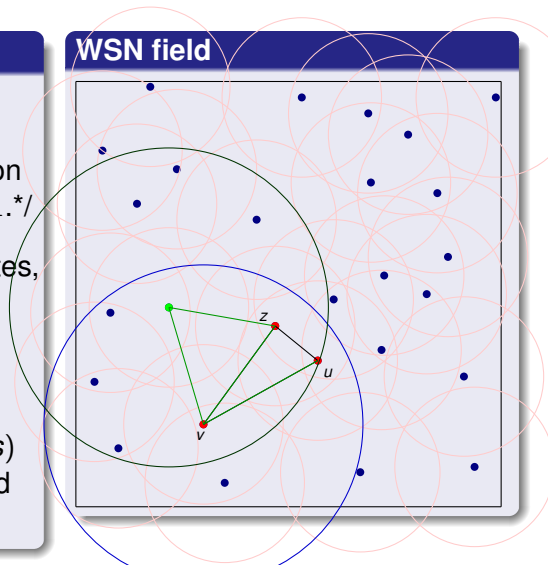
Mark s_2 */** The mark denotes, s_2 is visited but not fixed.**/*

Set $p_3 \leftarrow \text{null}$

$p, q \leftarrow$ two positions of s_3
from $\Delta s_1 s_2 s_3$.

Say $s_4 \in C(s_1, 2s) \cap C(s_3, 2s)$
 $-\{s_1, s_2, s_3\}$ is unmarked

WSN field



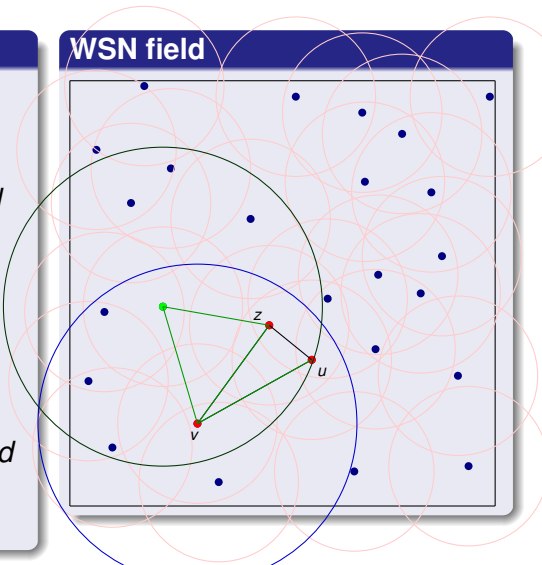
Localization under the Restricted Model

Localization Algorithm

Theorem

Any sensor at A , has at least three neighbours within $C(A, 2s)$, such that the neighbours can be connected to form a simple polygon around v_1 , with each side of length no more than $2s$. It is also possible to choose the neighbours such that the polygon encloses A . The polygon along with v_1 is a rigid subgraph (**rigid wheel** or simply **wheel**).

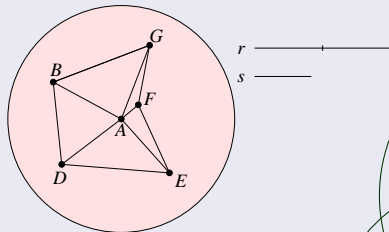
WSN field



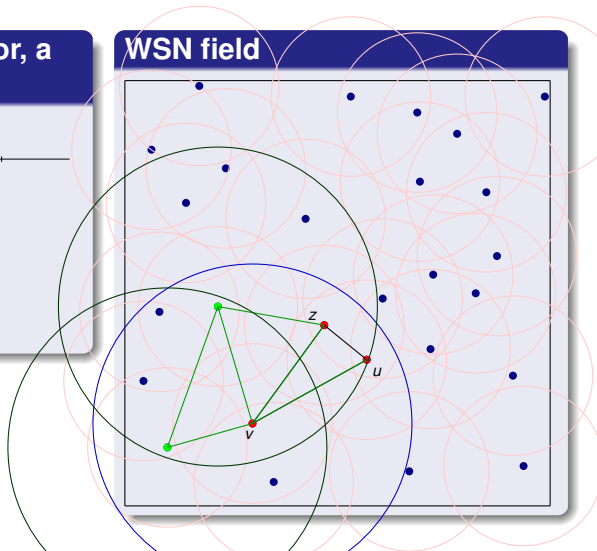
Localization under the Restricted Model

Localization Algorithm

Polygon around a sensor, a wheel.



WSN field



Localization under the Restricted Model

Localization Algorithm

```
func GETPOS( $s_3, s_1, s_2, p_2$ )
```

```
  if( $\Phi(s_4)$ ) /*  $\Phi(s_4)$  is fixed */
```

```
    if( $dist(p, \Phi(s_4)) = w(s_3, s_4)$ )
```

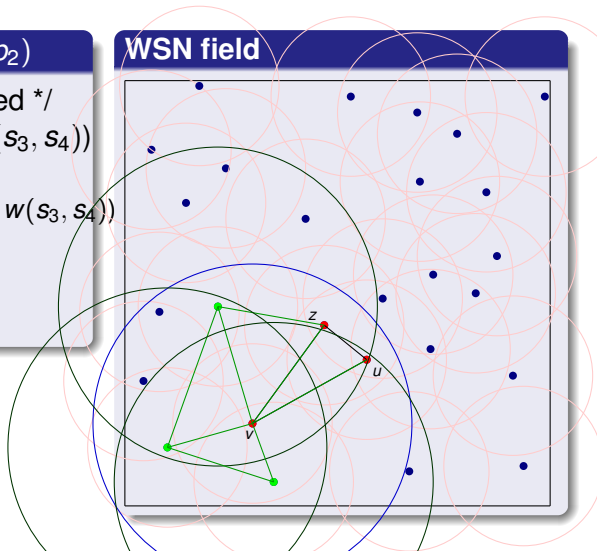
```
      Set  $p_3 \leftarrow p$ 
```

```
    else if( $dist(q, \Phi(s_4)) = w(s_3, s_4)$ )
```

```
      Set  $p_3 \leftarrow q$ 
```

```
    end if
```

WSN field



Localization under the Restricted Model

Localization Algorithm

```
func GETPOS( $s_3, s_1, s_2, p_2$ )
```

```
else /*  $\Phi(s_4)$  is not fixed */
```

```
  if( $t \leftarrow$  GETPOS( $s_4, s_1, s_3, p$ ))
```

```
    Set  $\Phi(s_4) \leftarrow t$ 
```

```
    Push  $s_4$  to located and locateIn2s
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    Push  $s_4$  to located and locateIn2s
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```
    Set  $p_3 \leftarrow q$ 
```

```
  end if
```

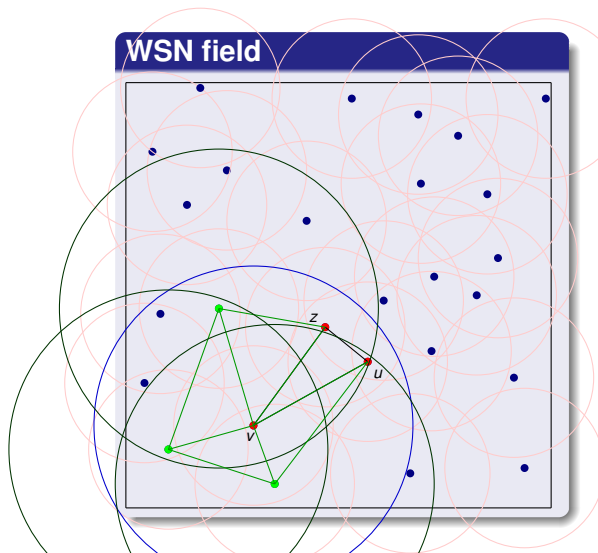
```
end if
```

```
Unmark  $s_2$ 
```

```
return  $p_3$ 
```

Localization under the Restricted Model

Localization Algorithm



Outline

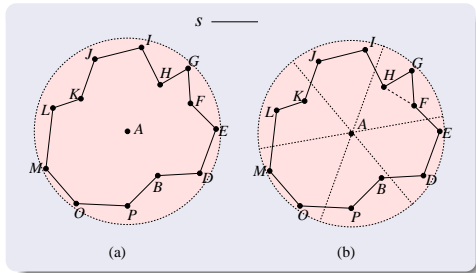
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Restricted Model

Correctness and progress

Theorem

Under the restricted model, each sensor will be surrounded by at least one rigid wheel consisting of no more than twelve vertices in the boundary.

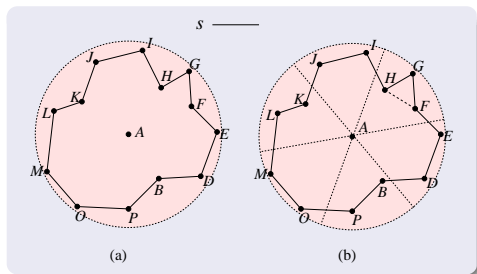


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Termination

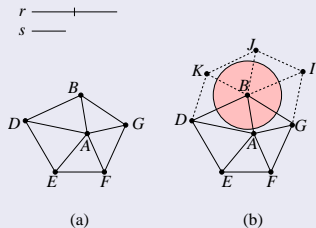
- The recursion $\text{GETPOS}(s_3, s_1, s_2, p_2)$ terminates after 12 calls at most.

Restricted Model

Correctness and progress

Theorem

At any point of time, if the network is not completely localized, there is at least one sensor which can be localized.

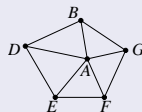
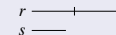


Restricted Model

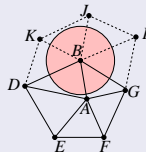
Correctness and progress

Theorem

At any point of time, if the network is not completely localized, there is at least one sensor which can be localized.



(a)



(b)

Progress

- at least one sensor with unknown position will be localized in each twelve cycles

Restricted Model

Complexity of the Algorithm

Finding Positions

```

while(located  $\neq$   $\emptyset$ )
  Push Top(located) to locateIn2s
  while(locateIn2s  $\neq$   $\emptyset$ )
     $s_2 \leftarrow$  Pop(locateIn2s)
    for all ( $s_3 \in C(s_1, 2s) \cap C(s_2, 2s)$ 
      and  $\Phi(s_3) = \text{null}$ )
       $\Phi(s_3) \leftarrow$  GETPOS( $s_3, s_1, s_2, \Phi(s_2)$ )
      Push  $s_3$  to located
      and to locateIn2s
    end for
  end while
   $s_1 \leftarrow$  Pop(located)
end while
  
```

- 1 The **for all** loop block will be executed in total (considering all repetitions of the earlier **while** loops) $|V|$ times, once for each vertex.
- 2 Checking in **for all** loop may have to be carried out is $O(|E|)$.

Restricted Model

Complexity of the Algorithm

Over all run time

Hence the worst-case time complexity of the algorithm is $O(|E|)$.

Outline

- 1 Introduction
- 2 Basic Model and Problem Statement
- 3 Localization under the Basic Model
- 4 Restricted Model**
 - Localization under the Restricted Model
 - Correctness and Complexity
 - **Simulation Studies**
- 5 Conclusion

Restricted Model

Simulation Studies

- 1 the worst case complexity, i.e., $O(n^2)$ may be achieved if the WSN has $O(n)$ sensors each with $O(n)$ neighbours.
- 2 Under random uniform deployment of nodes, the expected value of $|E|$ will be $O(n^2 \frac{r^2}{R})$, R is the field area.
- 3 For fixed r , the expected value of $|E|$ much lower than $O(n^2)$, in large scale WSN.

Restricted Model

Simulation Studies

Simulation Environment

- 1 We deploy sensors randomly with uniform distribution till the field is fully covered.
- 2 We repeat the experiment 30,000 times.
- 3 We carried out simulation studies (using a C++ program).

Simulation parameters

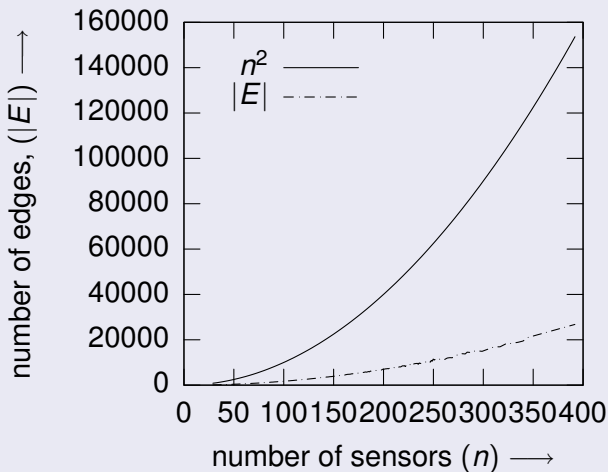
With the parameters :

Parameter	Value
WSN Field dimension	$50 \times 50 \text{ unit}^2$
Communication range of a sensor, r	24 unit
Sensing range of a sensor, s	10 unit
Range ratio, $\frac{r}{s}$	2.4

Restricted Model

Simulation Studies

Number of sensors vs. number of edges with length $\leq 2s$, where $s = 10$ unit, $R = 2500$ unit²



Restricted Model

Simulation Studies

Observation

$|E|$ is much lower than n^2 .

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Conclusion

Current Model

- The WSN is **fully covered**.
- Distances are measured exactly.

Conclusion

Current Model

- The WSN is **fully covered**.
- Distances are measured exactly.

Future Work

- The WSN has small holes.
- Distances have errors.

Thank you !